Preorder traversal:

Visit the root first, then all subtrees (from left to right) in preorder.

Postorder traversal:

Visit first all subtrees (from left to right) in postorder, then visit roots.

Inorder traversal (binary trees only):

Visit the left subtree, then the root, then the right subtree.

Polish notation corresponds a preorder traversal of the binary tree representing an algebraic expression. Standard notation (with brackets) corresponds to inorder traversal.

## Depth-first Search

Input: A connected graph G = (V, E),  $V = \{v_1, \ldots, v_n\}$ .

OUTPUT: A (depth-first) spanning tree T of G.

- (Step 1) Initialize  $T = (V_T, E_T)$ :  $V_T = \{v_1\}$  and  $E_T = \emptyset$ . Set  $current := v_1$ .
- (Step 2) Select the neighbour of current with the lowest index which is not yet visited. If no such neighbour exist, continue to Step 3. If it does exist, say  $v_i$ , set current :=  $v_i$  and return to Step 2.
- (Step 3) If  $current \neq v_1$ , then find the parent  $v_j$  of current, and set  $current = v_j$  (backtracking). Go to Step 2.
- (Step 4) If  $current = v_1$ , then stop.

## **Breadth-first Search**

INPUT: A connected graph G = (V, E),  $V = \{v_1, \ldots, v_n\}$ .

Output: A (breadth-first) spanning tree T of G.

- (Step 1) Initialize  $T = (V_T, E_T)$ :  $V_T = \{v_1\}$  and  $E_T = \emptyset$ . Insert  $v_1$  in a queue Q.
- (Step 2) Set current to be the vertex at the front of the queue Q, and delete this vertex from Q. Add all neighbours of current that have not been visited to the rear of the queue, in order of increasing index. If no such neighbours exist, repeat Step 2.

(Step 3) If Q is empty, then stop.

**Vertex set:**  $\{a, b, c, d, e, f, g, h\}$ .

Adjacency matrix:

 $\begin{bmatrix} 0 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 & 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 1 & 1 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 \end{bmatrix}$