# Depth-first Search

INPUT: A connected graph G = (V, E),  $V = \{v_1, \ldots, v_n\}$ .

OUTPUT: A (depth-first) spanning tree T of G.

- (Step 1) Initialize  $T = (V_T, E_T)$ :  $V_T = \{v_1\}$  and  $E_T = \emptyset$ . Set  $current := v_1$ .
- (Step 2) Select the neighbour of current with the lowest index which is not yet visited. If no such neighbour exist, continue to Step 3. If it does exist, say  $v_i$ , set current :=  $v_i$  and return to Step 2.
- (Step 3) If  $current \neq v_1$ , then find the parent  $v_j$  of current, and set  $current = v_j$  (backtracking). Go to Step 2.
- (Step 4) If  $current = v_1$ , then stop.

### **Breadth-first Search**

INPUT: A connected graph G = (V, E),  $V = \{v_1, \ldots, v_n\}$ .

Output: A (breadth-first) spanning tree T of G.

- (Step 1) Initialize  $T = (V_T, E_T)$ :  $V_T = \{v_1\}$  and  $E_T = \emptyset$ . Insert  $v_1$  in a queue Q.
- (Step 2) Set current to be the vertex at the front of the queue Q, and delete this vertex from Q. Add all neighbours of current that have not been visited to the rear of the queue, in order of increasing index. If no such neighbours exist, repeat Step 2.

(Step 3) If Q is empty, then stop.

**Vertex set:**  $\{a, b, c, d, e, f, g, h\}$ .

Adjacency matrix:

 $\begin{bmatrix} 0 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 & 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 1 & 1 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 \end{bmatrix}$ 

## Merge sort:

ished.

INPUT: A list L of n unordered numbers Output: A list OL of the same numbers in increasing order.

- (Step 1:) Form a binary tree T of height t, where  $t = \lceil \log_2(n) \rceil$ . Label its leaves with the elements of L. Label all intermediate nodes as unfinished, all leaves as finished.
- (Step 2:) Find an unfinished intermediate node v with two finished children with lists  $L_1$  and  $L_2$ :
  - (Step2a:) Merge  $L_1$  and  $L_2$  into L: take the smallest of the first elements of  $L_1$  and  $L_2$ , remove it from its list and put it at the end of L. Repeate this until  $L_1$  and  $L_2$  are empty.
  - (Step 2b:) Assign L to v, and label v as finished. Repeat this step until all nodes are fin-

(Step 3:) Let OL be the list assigned to the root.

A relation from set A to set B is a subset of  $A \times B$ .

Notation: Given a relation  $R \subseteq A \times B$ ,

$$xRy \Leftrightarrow (x,y) \in R$$

A relation R on a set A is a relation from A to A. Such a relation can have the following properties (universe of the quantifiers is A):

Reflexive:

$$\forall_x (x, x) \in R.$$

Symmetric:

$$\forall_x \forall_y (x, y) \in R \to (y, x) \in R.$$

**Antisymmetric:** 

$$\forall_x \forall_y [(x,y) \in R \land (y,x) \in R] \to x = y$$

Transitive:

$$\forall_x \forall_y \forall_z [(x,y) \in R \land (y,z) \in R] \to (x,z) \in R$$

A	xRy if:	refl.	sym.	antis.	tr
people	x is taller than y			X	
reals	x < y			X	
reals	$x \ge y$	X		X	
integers	x y	X		X	
integers	x + y is even	X	X		
people	x has the same age as $y$	X		X	
people	x has a parent in common with $y$	X		X	
integers	$y = x^2$				
sets	$x \subseteq y$	X		X	
integers	x and $y$ are				
	relatively prime		$\mathbf{X}$		
integers	x and $y$ , divided by 5	X	X		
	have the same remainder				
bit strings	x and $y$ have the same	X	X		
of length 7	number of 1's				
web pages	x has a hyperlink to $y$				

Partial order: a relation that is reflexive, antisymmetric and transitive.

#### **Examples:**

- The relation R on  $\mathbb{Z}$  where xRy if  $x \geq y$ ,
- The relation R on the collection of alls sets where xRy if  $x \subseteq y$ ,
- The relation R on  $\mathbb{Z}$  where xRy if x|y.

Equivalence relation: a relation that is reflexive, symmetric and transitive.

### **Examples:**

- The relation R on  $\mathbb{Z}$  where xRy if x-y is divisible by 5.
- The relation R on the set of all people where xRy if x and y have the same age.

A relation R on a set A can be represented as a directed graph  $G_R = (A, E)$  as follows:

The vertices of  $G_R$  are the elements of AThe edge set  $E = \{(a, b) \in A \times A | aRb\}$ .

A relation R on a set  $A = \{a_1, a_2, \dots, a_n\}$  can be represented by an  $n \times n$  relation matrix A as follows:

$$A_{i,j} = \begin{cases} 1 & \text{if } a_i R a_j \\ 0 & \text{otherwise} \end{cases}$$

Note that the relation matrix of R is the adjacency matrix of the directed graph representing R.