The Average Order of Dominating Sets of a Graph

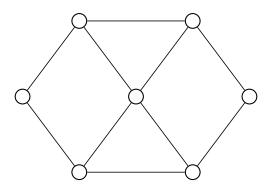
lain Beaton

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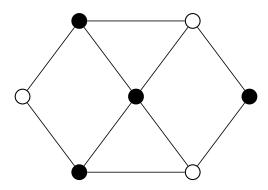
Atlantic Graph Theory Seminar, October 28, 2020

Joint work with Jason I. Brown

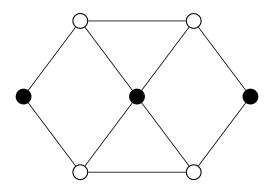
Dominating Set



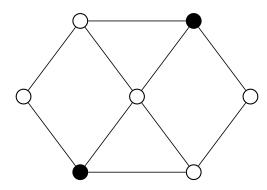
Dominating Sets



Dominating Sets

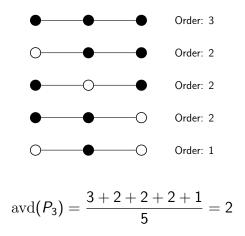


Dominating Sets



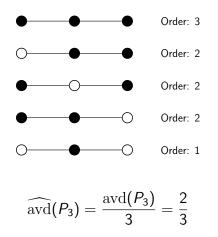
$\operatorname{avd}(G)$

The average order of the dominating sets of a graph G, denoted $\operatorname{avd}(G)$, is simply the sum of the orders of each dominating set divided by the number of dominating sets in G.





The normalized average order of the dominating sets of a graph G, denoted $\widehat{\text{avd}}(G)$, is simply $\operatorname{avd}(G)$ divided by the number of vertices in G.



The Domination Polynomial

Let d_k denote the number of dominating sets of order k in a graph G with n vertices. Then

$$\operatorname{avd}(G) = rac{\sum\limits_{k=0}^{n} k \cdot d_k}{\sum\limits_{k=0}^{n} d_k}$$

The **domination polynomial** D(G, x) of G is defined as

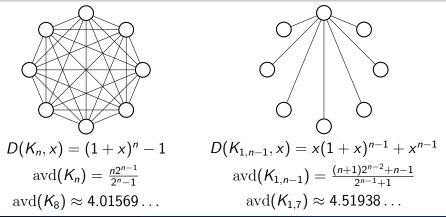
$$D(G,x) = \sum_{k=0}^{n} d_k x^k \Rightarrow \operatorname{avd}(G) = \frac{D'(G,1)}{D(G,1)}$$

avd(G) and the Domination Polynomial

For graphs G and H, $D(G \cup H, x) = D(G, x) \cdot D(H, x)$.

Lemma (B., Brown, 2020+)

Let G and H be graphs. Then $\operatorname{avd}(G \cup H) = \operatorname{avd}(G) + \operatorname{avd}(H)$.



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Other Average Graph Parameters

There has been work done on the averages of several other graphs invariants:

- The mean distance between vertices in a graph was introduced in 1977 by Doyle and Graver.
- The mean subtree order of a graph was introduced in 1983 by Jamison.
- The average size of an independent set in a graph was introduced in 2019 by Andriantiana et. el..
- Andriantiana et. el. also introduced the average size of a matching in a graph In 2020.

Other Average Graph Parameters

A common question is for each order n, which graphs maximize and minimize the given average graph parameter.

Average	Family	Maximum	Minimum	Author
Parameter		Graph	Graph	
Distance	Connected	P _n	K _n	Doyle et. el.
	graphs			(1977)
Subtree	Trees	???	P _n	Jamison
order				(1983)
Independent				Andriantiana
sets and	All graphs	$\overline{K_n}$	K _n	et. el.
Matchings				(2019,2020)
Independent				Andriantiana
sets and	Trees	P_n	$K_{1,n-1}$	et. el.
Matchings				(2019,2020)

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Extremal Graphs

For a graph G of order n, the upper bound for avd(G) is fairly straightforward:

$$\operatorname{avd}(G) \leq n = \operatorname{avd}(\overline{K_n}).$$

The lower bound for avd(G) is as you would expect:

Theorem (B., Brown, 2020+)

Let G be a graph of order n then $\operatorname{avd}(G) \ge \frac{n2^{n-1}}{2^n-1}$ with equality if and only if $G \cong K_n$.

Sketch of Proof

Proposition (B., Brown, 2020+)

For a graph G of order n, $d_k \leq d_{n-k}$ for all $k \leq \frac{n}{2}$.

Intuitively:

$$d_0, d_1, \dots, \underbrace{d_k, \dots, d_{\frac{n}{2}-1}, d_{\frac{n}{2}}, d_{\frac{n}{2}+1}, \dots, d_{n-k}, \dots, d_{n-1}, d_n}_{d_k \leq d_{n-k}}$$

Algebraically it means $k \cdot d_k + (n-k) \cdot d_{n-k} \ge \frac{n}{2}(d_k + d_{n-k})$, Therefore

$$\operatorname{avd}(G) = \frac{\sum\limits_{k=0}^{n} k \cdot d_{k}}{\sum\limits_{k=0}^{n} d_{k}} \ge \frac{\sum\limits_{k=0}^{n} \frac{n}{2} \cdot d_{k}}{\sum\limits_{k=0}^{n} d_{k}} = \frac{n}{2}$$

Extremal Trees

A similar, but more involved approach works to find the tree T on n vertices which minimizes $\operatorname{avd}(T)$.

Lemma (B., Brown, 2020+)

If T is a tree with n vertices then $d_{n-k} \ge d_{k+1}$ for all $k+1 \le \frac{n+1}{2}$.

Theorem (B., Brown, 2020+)

If T is a tree with n vertices $\operatorname{avd}(T) \ge \operatorname{avd}(K_{1,n-1})$ with equality if and only if $T \cong K_{1,n-1}$.

But which tree maximizes $\operatorname{avd}(T)$?

A natural guess would be paths...

$\operatorname{avd}(P_n)$

Theorem (Alikhani, Peng, 2009)

For all
$$n \ge 4$$
, $D(P_n, x) = x(D(P_{n-1}, x) + D(P_{n-2}, x) + D(P_{n-3}, x)).$

The general solution for the recurrence $D(P_n, x)$ satisfies is: $\lambda_1(x) = \frac{x}{3} + p(x) + q(x),$

$$\lambda_2(x) = \frac{x}{3} - p(x) - q(x) + \frac{\sqrt{3}}{2} (p(x) - q(x)) i,$$

$$\lambda_3(x) = \frac{x}{3} - p(x) - q(x) - \frac{\sqrt{3}}{2} (p(x) - q(x)) i,$$

where

$$p(x) = \sqrt[3]{\frac{x^3}{27} + \frac{x^2}{6} + \frac{x}{2} + \sqrt{\frac{x^4}{36} + \frac{7x^3}{54} + \frac{x^2}{4}}}$$
$$q(x) = \sqrt[3]{\frac{x^3}{27} + \frac{x^2}{6} + \frac{x}{2} - \sqrt{\frac{x^4}{36} + \frac{7x^3}{54} + \frac{x^2}{4}}}$$

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Theorem (B., Brown, 2020+)

Suppose functions $f_n(x)$ satisfy

$$f_n(x) = lpha_1(x)(\lambda_1(x))^n + lpha_2(x)(\lambda_2(x))^n + \cdots + lpha_k(x)(\lambda_k(x))^n$$

where $\alpha_i(x)$ and $\lambda_i(x)$ are fixed non-zero analytic functions, such that $|\lambda_1(1)| > |\lambda_i(1)|$ for all i > 1. Then

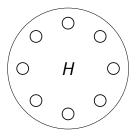
$$\lim_{n\to\infty}\frac{f_n'(1)}{nf_n(1)}=\frac{\lambda_1'(1)}{\lambda_1(1)}.$$

Theorem (B., Brown, 2020+)

$$\lim_{n\to\infty}\frac{\operatorname{avd}(P_n)}{n}=\lim_{n\to\infty}\widehat{\operatorname{avd}}(P_n)\approx 0.618419922.$$

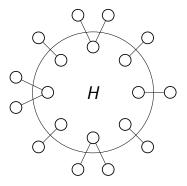
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A graph is called **sunlike** if it is formed by adding one or two pendant vertices to every vertex some base graph H.



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A graph is called **sunlike** if it is formed by adding one or two pendant vertices to every vertex some base graph H.



All sunlike graphs G have $\widehat{\operatorname{avd}}(G) = \frac{2}{3}$.

It remains an open question if there is a tree T with $avd(T) > \frac{2}{3}$.

Bounds

We know $\frac{n}{2} < \operatorname{avd}(G) \le n$, but can we do better?

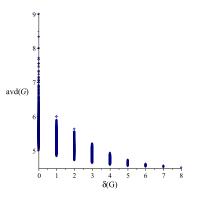


Figure 1: $\operatorname{avd}(G)$ over the minimum degree $\delta(G)$ for graphs of order 9.

Conjecture

Let G be a graph with $n \ge 2$ vertices. If G has no isolated vertices (so, in particular, if G is connected) then $\operatorname{avd}(G) \le \frac{2n}{3}$.

Lets take a little journey...

For a dominating set S, let a(S) denote the **domination critical** vertices of S. That is

 $a(S) = \{v \in S : S - v \text{ is not a dominating set}\}.$ 5 6 3 If $S = \{1, 5, 7\}$ then $a(S) = \{1, 7\}$

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For a graph G, let $\mathcal{D}_k(G)$ denote the collection of dominating sets of order k.

$$\sum_{S\in\mathcal{D}_k(G)}|\mathsf{a}(S)|=k\cdot d_k-(n-k+1)\cdot d_{k-1}.$$

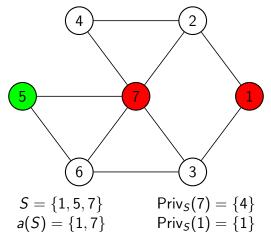
Lemma (B., Brown, 2020+)

For a graph G with n vertices.

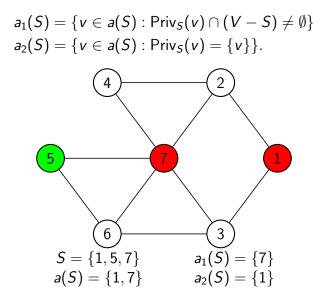
$$\sum_{S\in\mathcal{D}(G)}|a(S)|=2D'(G,1)-nD(G,1).$$

Where $\mathcal{D}(G)$ denotes the collection of all dominating sets of G

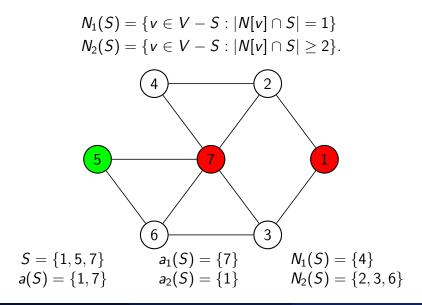
For a dominating set S and $v \in a(S)$ there is a non-empty set of vertices which are not dominated by S - v. We call these the **private neighbours** of v in S and denote them by $Priv_S(v)$.



We can then partition a(S) into two parts:



Furthermore, we can then partition V - S into two parts:



Intuition for $a_1(S)$ and $N_1(S)$

For a dominating set *S*:

- $a_1(S)$: Vertices in S which have private neighbours outside S.
- $N_1(S)$: Vertices not in S with one neighbours in S.
 - $N_1(S)$: Vertices which the are private neighbours outside of S.

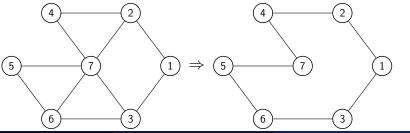
Lemma (B., Brown, 2020+)

Let G be a graph. For any dominating set, $|a_1(S)| \le |N_1(S)|$.

Lemma (B., Brown, 2020+)

Let G be a graph with is quasi-regularizable. For any dominating set, $|a_2(S)| \le |N_2(S)|$.

A graph G is called **quasi-regularizable** if one can replace each edge of G with a non-negative number of parallel copies, so as to obtain a regular multigraph of minimum degree at least one.



If G is quasi-regularizable then got every dominating set S:

$$|a(S)| = |a_1(S)| + |a_2(S)| \le |N_1(S)| + |N_2(S)| = n - |S|.$$

So if we sum this over all dominating sets be get

$$\sum_{S\in\mathcal{D}(G)}|\mathsf{a}(S)|\leq \sum_{S\in\mathcal{D}(G)}(n-|S|)=\sum_{k=0}^n(n-k)\cdot d_k=nD(G,1)-D'(G,1).$$

Then if we combine this with the previous result:

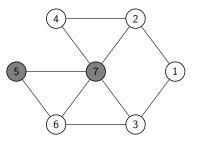
$$\sum_{S\in\mathcal{D}(G)}|a(S)|=2D'(G,1)-nD(G,1).$$

We obtain:

Theorem (B., Brown, 2020+) If G is a quasi-regularizable graph then $\operatorname{avd}(G) \leq \frac{2n}{3}$.

Bounding $a_2(S)$ for general graphs

Let $p_v(G)$ denote the collection of dominating sets of G - v which are not dominating sets of G. Below is an example of $p_1(G)$.



For each dominating set S and $v \in a_2(S)$, S - v is set which dominates G - v but not G.

Lemma (B., Brown, 2020+)

For any graph G be a graph, $\sum_{S \in \mathcal{D}(G)} |a_2(S)| = \sum_{v \in V(G)} |p_v(G)|$

Note that for each vertex v, $(2^{\deg(\nu)+1}-1)|p_{\nu}(G)| \leq D(G,1)$

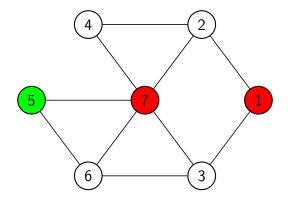
Theorem (B., Brown, 2020+)

Let G be a graph with $n \ge 2$ vertices and minimum degree $\delta \ge 1$. Then

$$\operatorname{avd}(G) \leq rac{2n(2^{\delta}-1)+n}{3(2^{\delta}-1)+1},$$

and so $\operatorname{avd}(G) \leq \frac{3n}{4}$.

We can do better...



 $\begin{array}{ll} S = \{1,5,7\} & a_1(S) = \{7\} & N_1(S) = \{4\} \\ a(S) = \{1,7\} & a_2(S) = \{1\} & N_2(S) = \{2,3,6\} \end{array}$

For any graph G,
$$\sum_{S \in \mathcal{D}(G)} |N_1(S)| = \sum_{e \in E(G)} |\mathcal{D}(G) - \mathcal{D}(G - e)|.$$

Lemma (Kotek, Preen, Simon, Tittmann, Trinks, 2012)

Let G be a graph. For every edge $e = \{u, v\}$ of G,

$$|\mathcal{D}(G)-\mathcal{D}(G-e)|=|p_u(G-e)|+|p_v(G-e)|-|p_u(G)|-|p_v(G)|.$$

Theorem (B., Brown, 2020+)

For any graph G with no isolated vertices,

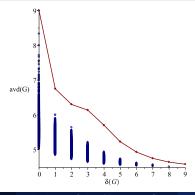
$$\operatorname{avd}(G) \leq \frac{n}{2} + \sum_{v \in V(G)} \frac{\operatorname{deg}(v)}{2^{\operatorname{deg}(v)+1} - 2}.$$

Corollary (B., Brown, 2020+)

For a graph G with minimum degree $\delta > 1$.

$$\operatorname{avd}(\mathcal{G}) \leq rac{n}{2}\left(1+rac{\delta}{2^{\delta}-1}
ight).$$

In particular if $\delta \geq 2\log_2(n)$ then $\operatorname{avd}(G) \leq \frac{n+1}{2}$.



Distribution of $\operatorname{avd}(G)$

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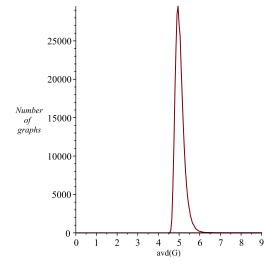


Figure 2: Distribution of avd(G) for all graphs of order 9

Let $\mathcal{G}(n, p)$ denote the sample space of random graphs on *n* vertices (each edge exists is independent present with probability *p*)

Theorem (B., Brown, 2020+)

For a fixed $p \in (0,1]$ let $G_n \in \mathcal{G}(n,p)$. Then

$$\lim_{n\to\infty}\widehat{\operatorname{avd}}(G_n)=\frac{1}{2}.$$

Proposition (B., Brown, 2020+)

The set
$$\left\{\widehat{\operatorname{avd}}(G) : G \text{ is a graph}\right\}$$
 is dense in $\left[\frac{1}{2}, 1\right]$.

Future Work

Future Work and Open Problems

- Can we show for all graphs with $\delta(G) \ge 0$, $\operatorname{avd}(G) \le \frac{2n}{3}$?
- Can we extend the work to the unimodality conjecture of the Domination Polynomial?

Conjecture (Alikhani, Peng, 2009)

The domination polynomial of any graph is unimodal.

$$\sum_{S\in\mathcal{D}_k(G)} |a(S)| = k\cdot d_k - (n-k+1)\cdot d_{k-1}.$$

• For a non-empty graph G, does there exists a vertex v and edge e such that

$$\operatorname{avd}(G - v) < \operatorname{avd}(G) < \operatorname{avd}(G - e).$$

THANK YOU!