Operational Theories of Physics as Categories

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The Plan

Find connections between:

- General probabilistic theories
- Categorical approaches, particularly effectus theory.

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Causality	Discard maps / Terminal object

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- ► Tests {f_x: A → B_x}_{x∈X}, X = finite outcome set Call subsets {f_y}_{y∈Y} ⊆ {f_x}_{x∈X} partial tests.
- ► Coarse-graining: Partial 'addition' $f \odot g : A \to B$ on events. $\{f_x : A \to B\}_{x \in X} \cup \{g_y\}_{y \in Y} \implies \{\bigcup_{x \in X} f_x\} \cup \{g_y\}_{y \in Y}$

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Examples

Many! Classical: deterministic or probabilistic. Quantum: Hilbert spaces or C*-algebras and c.p. sub-unital maps.

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Categorically, direct sums are finite coproducts (+, 0):

$$B_i \xrightarrow{\kappa_i} B_1 + \ldots + B_n = \bigoplus_{k=1}^n B_k \xrightarrow{\rhd_j} B_j , \quad \bowtie_j \circ \kappa_i = \begin{cases} \text{id} & i=j \\ 0 & i \neq j \end{cases}$$

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- ▶ \triangleright_i : $A_1 + ... + A_n \rightarrow A_i$ jointly monic
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- \blacktriangleright \otimes distributes over +

$$\blacktriangleright \doteqdot_{A+B} = [\doteqdot_A, \doteqdot_B], \doteqdot_I = \mathrm{id}, \doteqdot_{A\otimes B} = \lambda \circ (\doteqdot_A \otimes \doteqdot_B)$$

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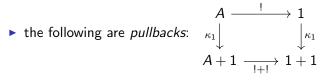
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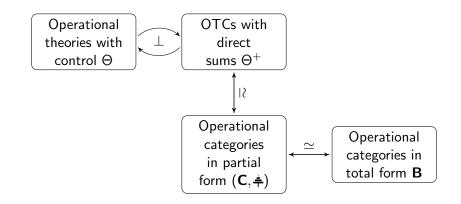
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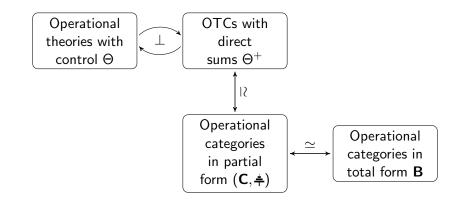
► the following are *pullbacks*: $A \xrightarrow{!} 1$ $\kappa_1 \downarrow \qquad \qquad \downarrow \kappa_1$ $A + 1 \xrightarrow{!+!} 1 + 1$

Weakening of notion of monoidal effectus (Jacobs et al.).

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Examples

Θ	В	С
Classical	Set	PFun
Quantum	$CStar^{\mathrm{op}}_{cpu}$	$CStar^{\mathrm{op}}_{cpsu}$

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