

Indefinite causal structures using diagrammatic methods

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The setting

- They own a lab,

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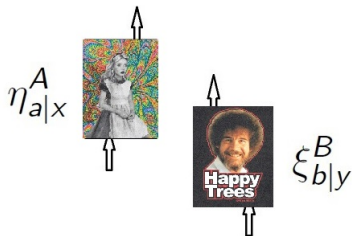
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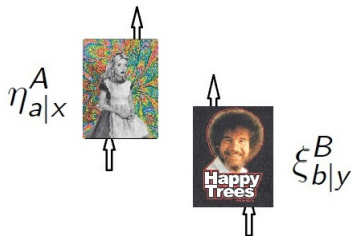
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Note:

We do not assume a fixed causal order between Alice and Bob. Only “local quantum mechanics”.

- Processes

$$\eta_{a|x}^A : B(H_{A_I}) \rightarrow B(H_{A_O})$$

$$\xi_{b|y}^B : B(H_{B_I}) \rightarrow (H_{B_O})$$

- Completely positive trace non-increasing

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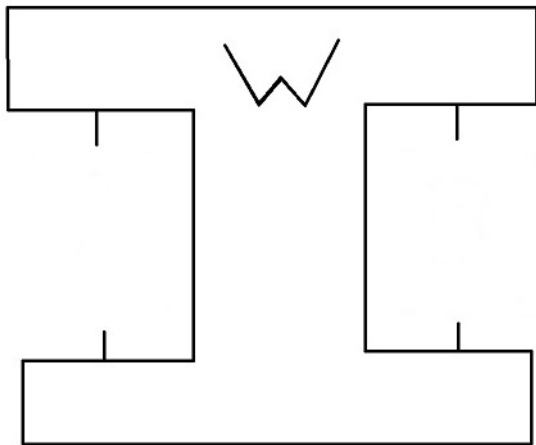
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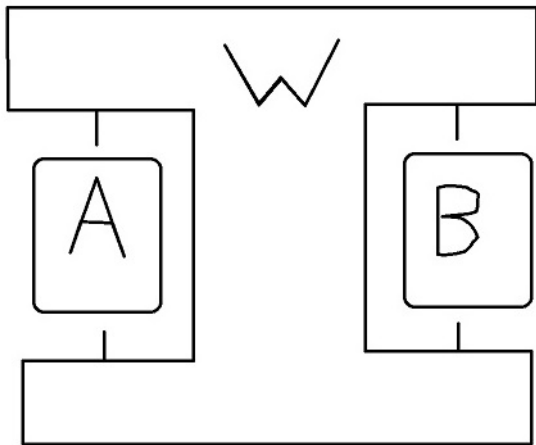
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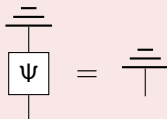
- W can be seen as a “generalized state”





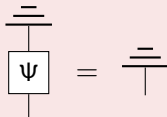
Definition

A process $\Psi : X \rightarrow Y$ is called *first order causal* if it is trace preserving. That is:



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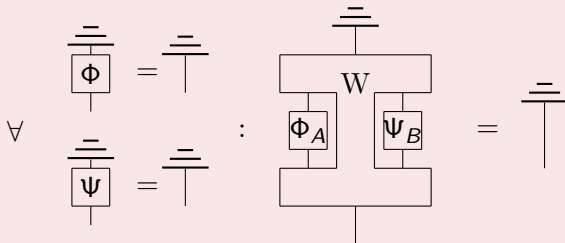
A process $\Psi : X \rightarrow Y$ is called *first order causal* if it is trace preserving. That is:



"If we discard the output of a causal process we might as well have discarded the input right away"

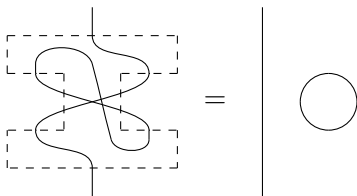
Definition

A process $W : (A_I \otimes B_I \rightarrow A_O \otimes B_O) \rightarrow (C_I \rightarrow C_O)$ is called *bipartite second-order causal (SOC₂)* if it sends causal processes to causal processes



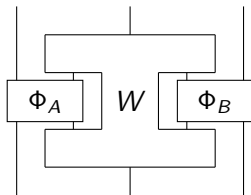
- Note SOC_2 only has to hold for separable processes

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- Example: swap

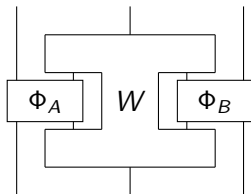


- (This implies it holds for non-signalling processes)

- What if A and B have ancilla systems



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Theorem

$$SOC_2 \Rightarrow \text{Completely } SOC_2$$

- We say a process theory has *enough causal states* if

$$\left(\forall \rho \text{ causal} \cdot \begin{array}{c} \square \Phi \\ \downarrow \\ \triangle \rho \end{array} = \begin{array}{c} \square \Phi' \\ \downarrow \\ \triangle \rho \end{array} \right) \implies \Phi = \Phi'$$

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- This is the case for operator algebras with cp maps

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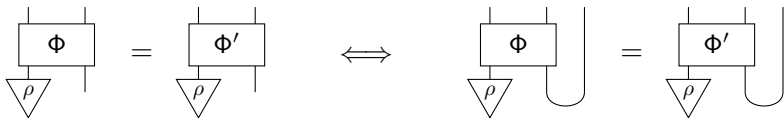
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Compactness gives: enough causal states \implies enough separable causal states

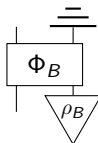
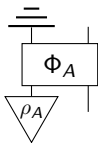
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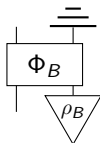
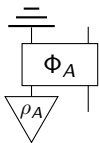


- For causal states ρ_A, ρ_B



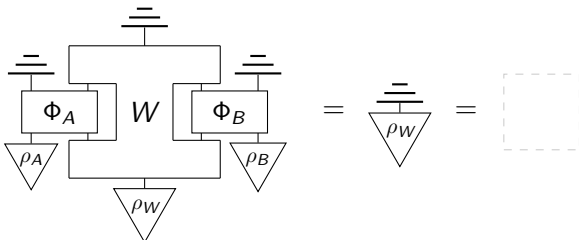
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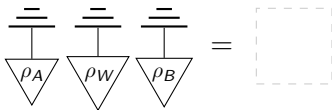


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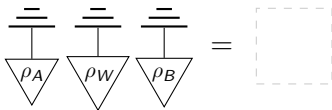
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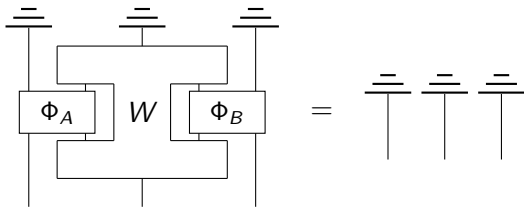
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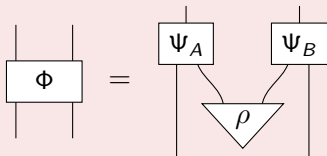


- Hence (enough separable causal states)



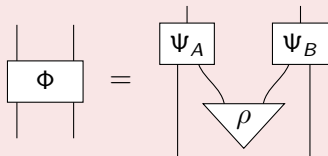
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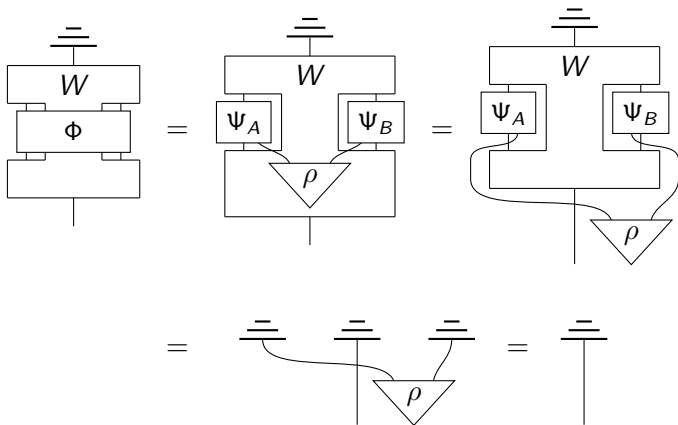
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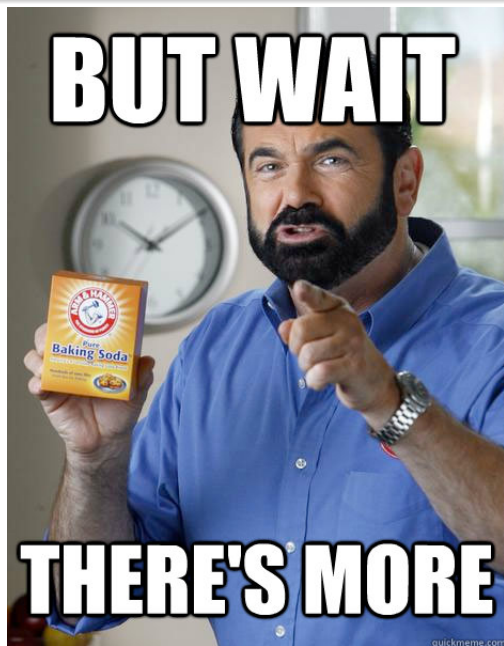
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I.e., Alice and Bob share a state, but no other communication.

If Φ is strongly non-signalling and W is SOC_2 then

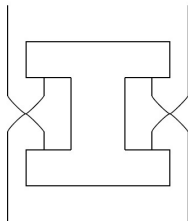




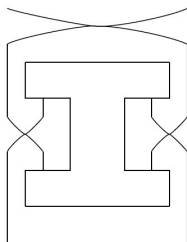
quickmeme.com



Swap maps are causal.
Plug them into W .



Then this is causal again.
And in fact non-signalling!



Bijection between W s and certain non-signaling boxes